

ADAIO



IPLS
INTERACTIVE PALLET LOADING SYSTEM

Research Report No. 81-9

by

Thom J. Hodgson

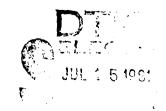
June, 1981

Department of Industrial and Systems Engineering
University of Florida
Gainesville, Florida 32611

APPROVED FOR PUBLIC RELEASE; DISTRIBUTION UNLIMITED

This research was supported in part by the U.S. Air Force, under contract number F73AFL-00360001, and by the Office of Naval Research, under contract number N00014-76-C-0096.

THE FINDINGS OF THIS REPORT ARE NOT TO BE CONSTRUED AS AN OFFICIAL DEPARTMENT OF THE AIR FORCE OR NAVY POSITION, UNLESS SO DESIGNATED BY OTHER AUTHORIZED DOCUMENTS.



9) R. son it repto

SECURITY CLASSIFICATION OF THIS PAGE (When Date Entered) READ INSTRUCTIONS REPORT DOCUMENTATION PAGE BEFORE COMPLETING FORM REPORT NUMBER 2. JOVT ACCESSION NO. J. RECIPIENT'S CATALOG NUMBER 81-9 10141 S. TYPE OF REPORT & PERIOD COVERED 4. TITLE (and Subtitle) Technical IPLS: INTERACTIVE PALLET LOADING SYSTEM PERFORMING ONG. REPORT NUMBER 81-9 W CONTRACT OR GRANT NUMBER(a) AUTHOR(4) F73AFL-00360001 (Air Force) Thom J./Hodgson N00014-76-C-0096 (Navy) PROGRAM ELEMENT, PROJECT, TASK AREA & WORK UNIT NUMBERS PERFORMING ORGANIZATION NAME AND ADDRESS Industrial and Systems Engineering University of Florida Gainesville, Florida 32611 11. CONTROLLING OFFICE NAME AND ADDRESS REPORT DATE Office of Naval Research - A.F. Logistics Jun**es 19**81 Arlington, VA Mgmt. Center 13. NUMBER OF PAGES Gunter AFS, Alabama 14. MONITORING AGENCY NAME & ADDRESS(If different from Controlling Office) 18. SECURITY CLASS. (of this report) Unclassified 15. DECLASSIFICATION/DOWNGRADING N/A 16. DISTRIBUTION STATEMENT (of this Report) APPROVED FOR PUBLIC RELEASE: DISTRIBUTION UNLIMITED. 17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20, If different from Report) N/A 18 SUPPLEMENTARY NOTES 19. FEY WORDS (Continue on reverse side if necessary and identify by block number) Pallet Loading Dynamic Programming Heuristic Real-Time 20 ABSTRACT (Continue on reverse side if necessary and identify by block number) :-The USAF Interactive Pallet Loading System (IPLS) is a real-time system for the planning of loading pallets. The system used a combination of Dynamic Programming and Heuristics, along with user interaction to accomplish the pallet loading task. The system is designed to be interfaced with the Base Automated Mobility System (BAMS). Examples are given. 7. TCV.

DD 1 JAN 73 1473 EDITION OF 1 NOV 48 IS OBSOLETE

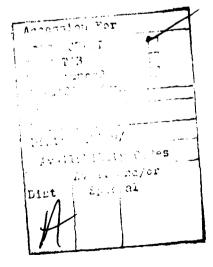
Unclassified

SECURITY CLASSIFICATION OF THIS PAGE (When Date Entered)

41 1311

TABLE OF CONTENTS

PA	GE.
ABSTRACT	i
INTRODUCTION	1
BACKGROUND	1
STRUCTURE AND USE OF IPLS	2
PALLET LOADING ALGORITHM	6
COMMENTS ON IMPLEMENTATION	1
REFERENCES	3



ABSTRACT

The USAF Interactive Pallet Loading System (IPLS) is a real-time system for the planning of loading pallets. The system used a combination of Dynamic Programming and Heuristics, along with user interaction to accomplish the pallet loading task. The system is designed to be interfaced with the Base Automated Mobility System (BAMS). Examples are given.

Introduction

The USAF Interactive Pallet Loading System (IFLS) is an interactive computerized system to assist in the Planning of IPLS is designed to interface with the USAF loading rallets. Base Automated Mobility System (BAMS) data base. The objective of IPLS is to automate the palletization of the unit material described in the BAMS data base, and to maximize the utilization of the available cube on the Fallet. IPLS is programmed in FORTRAN IV on a PDF-11/34. In what follows, previous efforts and current research on the mallet loading problem are discussed along with a seneral statement of the difficulty of the Problem and the philosophical approach taken here. In the following section the general structure and use of IPLS is presented. Then the mathematical structure of the pallet loading algorithm is laid out. Finally, comments on the requirements to implement IPLS in the "real world" are discussed along with rotential future enhancements. IFLS, at this point, should be considered a develormental system.

Background

The pallet loading problem is related to a problem long studied in the Operations Research literature: The Cutting Stock Problem. The cutting stock literature is not discussed here, but the interested reader is directed to a recent review paper by Golden [5]. The two-dimensional cutting problem (of which the pallet loading problem is a special case) has been studied by Christofides and Whitlock [1], Gilmore and Gomory [4], Hahn [6], and Herz [7]. The two-dimensional cutting problem becomes the pallet loading problem when the requirement for guillotine type cuts is dropped (i.e., straight cuts made in stages from one edge to the opposite edge of the object being cut, such as with a common paper cutter). Steudel [9] studied the pallet loading problem. However, his work was limited to the case where all boxes have the same length and width. He developed a procedure combining heuristics and dynamic programming.

DeSha [2], in an unpublished master's thesis, developed a heuristic for loading containers. His procedure first sets up stacks of items to fit the container height, then loads the stacks in the container to maximize the container floor area covered. The heuristic appears to obtain quite good results using a data base with boxes whose dimensions are randomly generated. The procedure does not, however, include considerations of center of gravity, positioning of hazardous material in the container, or box manipulation (no "this end up" assumption).

The pallet loading problem falls in the category of problems called NP-HARD [3]. Consequently, a truly efficient optimal place of the likely to be forthcoming. In developing an approach to the problem that would be consistent with the special needs of the USAF, it also became clear that it would be highly desirable for the system to be interactive. This would allow a user to guide the solution of a particular loading problem in order to deal with those unquantifiable clements of a real world loading problem. With these observations in mind,

Structure and Use of IPLS

IPLS is structured in two parts. The first part is the "BAMS Data Manipulator." The second in the "Pallet Loading Procedure." The Data Manipulator is a special purpose front end "sorter" that allows the user to obtain a set of candidate boxes for the Pallet Loading Procedure. The Data Manipulator is structured on the "menu" approach and uses a series of data bases in order to build the candidate set. The data bases are named BAMS.DAT, BAM1.DAT...., BAM4.DAT, and PACK.DAT. BAMS.DAT contains a card image file of the Base Automated Mobility System Data (obtained in the present case from Moody, AFB). BAM1.DAT is an unformatted copy of BAMS.DAT. BAM2.DAT...., BAM4.DAT are temporary files used in the selection of candidate boxes. PACK.DAT is used to store the BAMS data for boxes allocated to the pallet.

The menu (option list) of the Data Manipulator contains seven options:

1. IDENTIFY THE LIST OF ITEMS allows the user to choose which squadron increment(s) to include in the candidate data base. The load data for selected increment(s) are read from BAM1.DAT and written into BAM2.DAT; 2. SORT FILE OF BOXES

allows the user to sort the candidate data base by length-width or height (see figure 1a, and note that user responses are circled). This facilitates the choice process (particularly when coupled with option 4 in the menu). The sort is two phase (i.e., if the primary sort is by height, the secondary sort is by length-width within height, and visa-versa). A frequency analysis of the box dimensions is also provided on screen (figure 1a). The sort can be performed on BAM2.DAT,...,BAM4.DAT; 3. SELECT SUBSET OF BOXES FROM FILE

allows the user to select from BAM2.DAT (BAM3.DAT) and place the desired boxes in BAM3.DAT (BAM4.DAT). Selection is made by inputting a lower bound and upper bound on the desired dimension (length-width or height) for the selection process. An abbreviated frequency analysis of the box dimensions is provided on screen (figure 1b) to facilitate the choice process;

4. VIEW A "BAMS" FILE

allows the user to view the BAM1.DAT,...,BAM4.DAT, and PACK.DAT files (figures 2a and 2b);

5. MANIPULATE SPECIFIC BOX(ES) IN FILE

allows the user to eliminate box(es), to move box(es) from BAMX.DAT to BAMX+1.DAT, or change orientation of box(es) (figures 3a, 3b, and 3c);

5. STACK BOXES PRIOR TO FALLET LOADING

allows the user to create a file of "stacked" boxes (figure 4). The boxes are stacked, using a Dynamic Programming "knapsack" code [8], with the objective of maximizing the cube filled over a "base" box. The "stacks" can then be loaded using the Pallet Loading Procedure.

9. STOP

terminates the Data Manipulator.

- 1. IDENTIFY LIST OF ITEMS
- 2. SORT FILE OF BOXES
- 5. SELECT SUBSET OF BOXES FROM FILE
- 4. VIEW A "BAMS" FILE
- 5. MANIPULATE SPECIFIC BOX(ES) IN FILE
- 6. STACK BOXES PRIOR TO PALLET LOADING
- 9. STOP

(1/2/3/4/5/6/979(2)

MONT "BAM2", "BAM3", OR "BAM4" L2, 3, 40 ()

FRIMARY SORT BY "LIH-WTH", OR "HGT" EL,W,H) (L)

SECONDARY SORT BY "LTH-WTH", "HGT", OR NONE (L, W, H, P) (H)

LENGTH-WIDTH-HEIGHT ANALYSIS

1 , W . H	NO#	WEIGHT	AREA-FT	CUBE-FT	LyldaH	NO#	WEIGHT	AREA FT	<u>Eurr</u>
3	1	90.	11.	3.	3.2	2	170.	3.	• .
-4	3	125.	22.	7.	33	23	4659.	10:.	417
3	1.	50.	1.	1.	34	3	215.	·7.	÷ r
10	4	152.	34+	23.	35	10	1100.	21.	•••
11	5	382.	23.	21.	33	5	307.	17.	J
12	1.	89.	12.	12.	37	7	1125.	19.	59.
13	4	220,	6.	7.	30	.4	704.	34,	105.
14	1	130.	8.	10.	40	1	50.	٥.	i.
15	8	1019.	29.	36.	42	1	52.	2.	5.
16	14	1932.	55.	73.	43	33	3165.	75.	270.
17	27	2730.	114.	161.	46	4	1000.	11.	40.
18	45	4933.	241.	361.	47	ó	1200.	15.	59.
19	42	5178.	223.	353.	48	2	415.	9.	J6.
20	23	3067.	112.	187.	51	3	300.	12.	53.
21	11	1680.	50.	88.	52	1	130.	2.	10.
22	6	600.	80.	146.	56	2	760.	11.	52.
23	1	130.	5.	10.	5 <i>7</i>	1	267.	9.	41.
24	ර	535.	38.	76.	58	ర	600.	30.	146
25	4	625.	42.	86.	61	1	215.	4.	22.
26	1	75.	7.	15.	67	2	2120.	31.	174.
27	1	75.	3.	7.	68	2	2120.	31.	174.
28	1	400.	4.	10.	73	1	50.	1.	5.
29	1	50.	2.	5.	74	1	55.	2.	14.
30	2	190.	9.	24.	92	1	89.	2.	12.
31	12	1484.	27.	70.	O	0	0.	э.	0.
TOTAL	126	17431.		1388.					

Figure la

- 1. IDENTIFY LIST OF ITEMS
- 2. SORT FILE OF BOXES
- 3. SELECT SUBSET OF BOXES FROM FILE
- 4. VIEW A 'BAMS' FILE
- 5. MANIPULATE SPECIFIC BOX(ES) IN FILE
- 6. STACK BOXES PRIOR TO PALLET LOADERS
- 9. STOP

E1/2/3/4/5/6/917(3)

READ FROM "BAM2", WRITE TO "BAM3", OR READ FROM "BAM3", WRITE TO "BAM4" [2/3]

CHOOSE OPTION

- 1. SELECT ALL DATA
- 2. SELECT BY HEIGHT
- 3. SELECT BY LENGTH-WIDTH
- 9. RETURN

61,2,3,937 3

L,W,H	₩ 0#	1 y W y H	NO#	LyWyH	# 0 <i>N</i>	L. , W , H	NO#
10	4	18	18	2.1	2	43	33
17	11	17	29	36	2	()	0

CHOOSE LOWER BOUND (17)

CHOOSE UPPER BOUND (19

CHOOSE OPTION

- 1. IDENTIFY LIST OF ITEMS
- 2. SORT FILE OF BOXES
- 3. SELECT SUBSET OF BOXES FROM FILE
- 4. VIEW A "BAMS" FILE
- 5. MANIPULATE SPECIFIC BUX(ES) IN FILE
- 6. STACK BOXES PRIOR TO PALLET LOADING
- 9. STOP

T1/2/3/4/5/6/939**(** 9

TT1 -- STOP

Figure 1b

- 1. IDENTIFY LIST OF ITEMS
- 2. SORT FILE OF BOXES
- 3. SELECT SUBSET OF BOXES FORM FILE
- 4. VIEW A "BAMS" FILE
- 5. MANIPULATE SPECIFIC BOX(ES) IN FILE
- 6. STACK BOXES FRIOR TO PALLET LOADING
- 9. STOP

E172/3/4/5/6/91**?(** 4

T

WANT "BAM1", "BAM2", BAM3", "BAM4", OR "PACK" E1/2/3/4/733

INCR	ITEM	CUBE	L.TH	WTH	HGT	WEIGHT	11	STACE
1415	NR	FT.	IN.	1N.	IN.	LBS.	1	CODE
B001	006	8	17	19	43	200		0
B001	013	8	17	19	43	200	C	0
B001	800	8	17	19	43	125		()
F001	009	8	17	19	43	125		0
B001	010	8	17	19	4.3	1.25		0
B001	011	8	17	19	43	1.25		()
B001	012	8	17	19	43	100		0
B001	007	8	17	19	43	85		0
1041	002	9	18	19	43	95		Ü
1041	003	9	18	19	43	90		0
1041	004	9	18	19	43	95		O
1041	005	9	18	19	43	95		O
1041	006	9	18	19	43	95		()
1041	007	9	18	19	43	95		0
1041	008	9	18	19	43	95		O
1041	009	9	18	19	43	95		′)
1041	010	9	18	19	43	95		Ö
I041	011	9	18	19	43	95		O
1041	012	9	18	19	43	95		()
1041	013	9	18	19	43	95		0

INCR	ITEM	CURE	LTH	WTH	HGT	WEIGHT	Н	STACK
NR	NR	FT.	IN.	IN.	IN.	LBS.	Z.	CODE
1041	014	9	18	19	43	25		()
1041	015	9	18	19	43	95		Ò
1041	016	9	18	19	43	95		()
1041	017	φ	1.8	1.9	43	25		()
1041	018	9	18	19	43	95	•	0
1041	019	Ģ	18	19	43	6.7		Ó
1041	020	3	17	19	43	100		()
1041	021	8	1.7	1.9	43	100		()
1041	022	3	12	19	43	46		0

Figure 2a

- 1. IDENTIFY LIST OF ITEMS
- 2. SORT FILE OF BOXES
- 3. SELECT SUBSET OF BOXES FROM FILE
- 4. VIEW A "BAMS" FILE
- 5. MANIPULATE SPECIFIC BOX(ES) IN FILE
- 6. STACK BOXES PRIOR TO MALLET LUADING
- 9. STOP

T 1/2/3/4/5/6/939(4)

WANT "BAM1", "BAM2", BAM3", "BAM4", OR "PACK" E1/2/3/4/737

INCR	ITEM	CUBE	LTH	WTH	HGT	WEIGHT	14	STACK	LOC	LEVEL
NR	NR	FT.	IN.	IN.	IN.	LBS.	Z	CODE	I. W	14.
B001	006	8	17	19	43	200		0	3600	0
B001	013	8	17	19	43	200	С	0	5300	0
B001	008	8	17	19	43	125		0	7000	0
B001	009	8	17	19	43	125		()	8700	0
B001	010	8	17	1.9	43	125		()	19	0
B001	011	8	17	19	43	125		Ü	1719	0
B001	012	8	17	1.9	43	100		()	3419	0
B001	007	8	17	19	43	85		()	5119	O
I 0 4 1	002	9	18	19	43	95		0	()	9
I041	003	9	18	19	43	95		()	1800	0
IO41	004	9	18	19	43	95		()	6819	()
I041	005	9	18	19	43	95		0	8619	()
I U 41	006	9	18	19	43	95		ð	38	0
I041	007	9	18	19	43	95		()	1938	O
I041	008	9	18	19	43	95		O	3838	Ö
1041	009	9	13	19	4.3	95		\dot{O}	5738	0
I041	010	9	18	19	43	95		$\langle \cdot \rangle$	7638	Ç
1041	011	9	18	18	43	25		Q	56	Ö
1041	012	9	18	19	43	95		()	1856	O
1041	013	9	1.8	19	43	95		()	3656	0

INCR	ITEM	CUBE	LTH	WIH	HGT	WEIGHT	H	STACK	LOC	LEVEL
NR	NR	FT.	IN.	IN.	IN.	LBS.	Z	CODE	i W	IN.
1041	014	9	18	19	43	95		0	5456	0
IO41	015	9	18	19	43	95		Ö	7256	()

Figure 2b

- 1. IDENTIFY LIST OF ITEMS
- 2. SORT FILE OF BOXES
- 3. SELECT SUBSET OF BOXES FROM FILE
- 4. VIEW A "BAMS" FILE
- 5. MANIPULATE SPECIFIC BOX(ES) IN FILL
- 6. STACK BOXES PRIOR TO PALLET LOADING
- 9. STOP

III/2/3/4/5/8/919(5

CHOOSE OFTION

- 1. ELIMINATE A BOX
- 2. COPY A BOX FROM BAMW.DAT TO BAMWHI.DAT
- 3. ROTATE SINGLE BOX, LIH TO HOT OR WITH TO HOT
- 4. ROTATE BOXES FOR COMMON HELGHT
- 9. RETURN

E1/2/3/4/**9]**(1

MANIPULATE BAM2, BAM3, OR BAM4 82,3,43 (74

	INCR	ITEM	CUBE	LTH	WIH	HST	WEIGHT	14	STACK
	NR	NR	FT.	IN.	IN.	IN.	LBS.	Z	CODE
1.	B001	006	8	17	19	43	200		0
2.	B001	013	8	17	19	43	200	C	0
3.	B001	800	8	17	19	43	125		0
4.	B001	009	8	17	19	43	125		Ü
5.	B001	010	8	1.7	19	43	1.25		Ō
6.	B001	011	8	1.7	19	43	125		Ű
7.	B001	012	8	17	1.9	43	100		Θ
8.	B001	007	8	17	i 9	43	85		J
9.	1041	002	Ş	18	1.5	43	φ_{ij}		()
10.	1041	003	9	13	1.9	43	95		Ģ
11.	1041	004	9	18	19	43	95		ા
12.	1041	005	9	18	19	43	95		0
13.	1041	006	9	18	19	43	3.2		Ü
14.	IO41	007	9	13	19	4.3	95		0
15.	1041	008	9	18	19	43	95		Ü

ENTER NUMBER OF BOX TO BE ELIMINATED OR O TO CONTINUE :

1 BOXES ELIMINATED

28 BOXES REMAIN IN DAMA.DAT

ARE ELIMINATIONS CY/NIT

Figure 3a

- 1. ELIMINATE A BOX
- 2. COPY A BOX FROM BAM*.DAT TO BAM*+1.DAT
- 3. ROTATE SINGLE BOX, LTH TO HGT OR WITH TO HGT
- 4. ROTATE BOXES FOR COMMON HEIGHT
- 9. RETURN

E1/2/3/4/937(2)

COPY FROM "BAM2" TO "BAM3" OR COPY FROM "BAM3" TO "BAM4" 62/319

	INCR	ITEM	CUBE	LTH	HTW	HGT	WEIGHT	Н	STACK
	NR	NR	FT.	IN.	IN.	IN.	LBS.	Z	CODE
1.	B001	015	9	10	36	43	46		0
2.	B001	017	9	10	36	43	46		0
3.	B001	006	8	17	19	43	200		0
4.	B001	013	8	17	19	43	200	C	0
5.	B001	800	8	17	19	43	125		0
6.	B001	009	8	17	19	43	125		0
7.	B001	010	8	17	19	43	125		0
8.	B001	011	8	17	19	43	125		O
9.	B001	012	8	1.7	1.9	43	100		0
10.	B001	007	8	17	19	43	85		0
11.	EA02	012	5	10	21	43	30		0
12.	EB02	012	5	10	21	43	30		0
13.	I041	002	9	18	19	43	95		()
14.	I041	003	9	18	19	43	95		O
15.	I041	004	9	18	19	43	95		()

ENTER NUMBER OF BOX TO BE COPIED FROM "BAM3" TO "BAM4" OR O TO CONTINUE 4

INCR ITEM CUBE LTH WTH HGTWEIGHT H STACK LBS. NR NR FT. IN. IN. IN. Z CODE 4. B001 013 17 Ü 1.9 43 200

1 BOXES COPIED INTO "BAM4"

COPY ANOTHER BOX CYZNIT N

- 1. ELIMINATE A BOX
- 2. COPY A BOX FROM BAM*.DAT TO BAM*+1.DAT
- 3. ROTATE SINGLE BOX, LTH TO HGT OR WITH TO HGT
- 4. ROTATE BOXES FOR COMMON HEIGHT
- 9. RETURN

E1/2/3/4/93?

MANIPULATE BAM2, BAM3, OR BAM4 02,3,41 (?4

ENTER UPPER BOUND ON COMMON HEIGHT (19

ENTER LOWER BOUND ON COMMON HEIGHT (12

29 BOXES IN "BAM4" HAVE BEEN ROTATED

DO YOU WISH TO VIEW "BAM4" [Y/N] N

CHOOSE OFTION

- 1. ELIMINATE A BOX
- 2. COPY A BOX FROM BAM*.DAT TO BAM*+1.DAT
- 3. ROTATE SINGLE BOX, LTH TO HGT OR WITH TO HGT
- 4. ROTATE BOXES FOR COMMON HEIGHT
- 9. RETURN

E1/2/3/4/937(9)

CHOOSE OFTION

- 1. IDENTIFY LIST OF ITEMS
- 2. SORT FILE OF BOXES
- 3. SELECT SUBSET OF BOXES FROM FILE
- 4. VIEW A "BAMS" FILE

STOP

- 5. MANIPULATE SPECIFIC BOX(ES) IN FILE
- 6. STACK BOXES PRIOR TO PALLET LOADING
- 9. STOP

£1/2/3/4/5/6/93f

TT1

Figure 3c

- 1. IDENTIFY LIST OF ITEMS
- 2. SORT FILE OF BOXES
- 3. SELECT SUBSET OF BOXES FROM FILE
- 4. VIEW A "BAMS" FILE
- 5. MANIPULATE SPECIFIC BOX(ES) IN FILE
- 6. STACK BOXES PRIOR TO PALLET LOADING
- 9. STOP

_1_1/2/3/4/5/6/90f(

WHICH FILE IS TO BE STACKED [2/3]

	INCR	ITEM	CUBE	LTH	WTH	HGT	WE LOHE	11	STACK
	NR	NR	FT.	IN.	IN.	Ti,	LBS.	Z	come
1	B001	015	9	1.0	36	43	46		O
2	B001	017	9	10	36	43	46		0
3	B001	006	8	17	19	43	200		O
4	B001	013	8	17	19	43	200	C	0
5	B001	008	8	17	19	43	125		()
6	B001	009	8	1.7	19	43	125		0
7	B001	010	8	17	19	43	125		O
8	B001	011	8	17	19	43	125		O
9	B001	012	8	17	19	43	100		0
10	B001	007	8	17	19	43	85		0
11	EA02	012	5	10	21	43	30		0
12	EB02	012	5	1.0	21	43	30		C
13	I041	002	9	18	19	43	95		0
14	1041	003	9	18	19	43	95		0
15	TO41	004	9	18	19	43	95		0

WANT ONE OF THESE FOR BASE?
IF SO, INPUT NUMBER, OTHERWISE HIT "RETURN" (13

INPUT MAXIMUM PALLET HEIGHT (INCHES) (96)

							WEIGHT	
14	I041	003	9	18	19	4.3	95	()
13	I041	002	9	18	19	43	95	()

WANT TO PUT STACK ON BAM4.DAT EY/NJF(Y)

How means contion lists of the Hallet Loading Procedure contains five options:

Howe the oren to load the collet using the set of cardidate or ea. This represents a forward bask on the Pallot Loader. The Pollet toader is structured as a Dynamic Program [8] and colvectwo dimensional loading problem (i.e., loads one layer of the ea. or stacks of boxes, on the sallet) (figures 5,6); to be a Pallot FOAD SOLUTION

The user to fetch the solution generated by the BYNALAAN contine. This represents a backward mass on the mallet loading the Aliast of loaded boxes and their locations (x,x condinates) is provided. A plan view micture of the boxes loaded on the mallet is provided as an option. If gures 6.7...

3. PUT PACKED BOXES ON PACK.DAT ciloaded boxes from the RAMS.DAT currently being used and places them in the PACK.DAT file. The user can specify the level (height) that the boxes are to be

slaced on the mallet. This facilitates multiple masses of the packing routine while loading a mallet in "lawers";
4. CHUNCH LOOSE BOXES IN SOLUTION TOWARD ORIGIN
allows the user to "tighten" a mallet load solution by shifting ho we toward the ummer left—hand corner (mlan view) of the mallet as much as mossible. A mlan view micture of the boxes loaded on the mallet is provided as an option; and 9. STOP

terminates the Pallet Loading Procedure.

The following describes a way of using the irls system. Assume that the BAMS.DAT file contains a list of equipment appropriate to the planned mission. If only a partial list is to be used, the system "editor" can be used to delete specific items from the file. Then the "READIT" program (see appendix) can be used to load the PAM1.DAT file from BAMS.DAT.

One approach to loading (using the system) would be, first, to identify a set of candidate boxes for loading and then to load the mallet(s) with boxes from the candidate list.

The first step in identifying a set of candidate boxes is to use the IDENTIFY THE LIST OF ITEMS option in the Data Manigulator to select the squadron increment(s) to be loaded on the pallet(s). This results in a subset of the BAMS.DAT boxes being stored in the RAM2.DAT file. Next. the boxes should be sorted using the SORT THE FILE OF BOXES option. By examining the frequency analysis of box dimensions (provided as part of the sort, figure 1a), the user starts to get an overall feeling as to the characteristics of the set of candidate boxes. Then using the VIEW A 'BAMS' FILE option (figure 2a) the user can further winhance his knowledge of the (sorted) candidate set. At this roint the user can take several different approaches to pallet loading. The following outlines three possible approaches. 1. Attempt to select a set of boxes with a common height. This will allow the user to load a layer of boxes on the pallet (using the pallet loading procedure). Then he/she can select another set of boxes and load another layer until the maximum mallet height is reached (figure 8a).

98 BEES

Collinate and Line

THE MATTER A TIST OF LITTING

- of and fill OF BOXES
- AN SELECT SOBSET OF BOXES FROM FILE
- 4. UTTW A "BAMS" FILE
- DE MANIPULATE SPECIFIC BOX(ES) IN FILE
- * STACE BOXES PROOF TO PALLET LOADING
- 9. 31.W

- F1/2/3/4/5/6/91**(** 4

WARD "BOND", "BAM2", BAM3", "BAM4", OR "PACK" 5

TRUR	LTEM	CUBIL	LTH	WTH	HGT	WEIGHT	14	STACK
NR	NR	FT.	TN.	IN.	IN.	LBS.	Ζ	CODE
B001	012	9	1.0	36	43	46		0
ROOL	008	8	1.7	19	43	200		0
8001	013	8	17	19	43	200	C	()
1001	008	3	17	19	43	125		0
B001	0.02	8	1.7	19	43	125		0
B001	010	8	1.7	19	43	125		0
B001	011	8	17	19	43	125		0
B001	012	8	1.7	19	43	100		0
R001	007	8	1.7	19	43	85		0
EB02	012	5	10	21	43	30		0
1041	009	9	18	17	43	95		0
IO41	010	9	18	19	43	95		0
1041	011	9	18	19	43	95		0
IO41	012	9	18	19	43	95		0
IO41	013	9	18	19	43	95		0
1041	014	9	18	13	43	95		0
1041	015	9	13	19	43	95		0
IO41	016	9	18	19	43	95		0
1041	017	9	18	19	43	95		0
1041	018	9	18	19	43	95		0

TNCR	TTEM	CUBE	LTH	WTH	HGT	WEIGHT	Н	STACK
NR	NR	FT.	IN.	IN.	IN.	LBS.	Z	CODE
I041	019	9	18	19	43	67		0
1041	020	8	17	19	43	100		0
I041	021	8	17	19	43	100		0
1041	022	8	1.7	19	43	46		0
1041	014	9	18	19	43	95		0

Figure 5

RUN LOADIT WHICH FILE: BAM1? BAM2? BAM3? BAM4? E1/2/3/40 ZERO PACK.DAT? CY/NI Y INPUT LENGTH(60 INPUT WIDTH CHOOSE OPTION 1. LOAD PALLET 2. GET FALLET LOAD SOLUTION 3. PUT PACKED BOXES ON PACK.DAT 4. CRUNCH LOOSE BOXES IN SOLUTION TOWARD ORIGIN 9. STOP E1/2/3/4/91 PURE RECTANGLE CY/NI Y INPUT LENGTH(60 INPUT WIDTH CHOOSE OFTION 1. LOAD PALLET 2. GET PALLET LOAD SOLUTION 3. PUT PACKED BOXES ON MACKIDAT 4. CRUNCH LOOSE BOXES IN SOLUTION TOWARD ORIGIN 9. STOP E1/2/3/4/91 LOCATION BOX# LW LTH WITH 1. 36 36 10 21 10 3636 10 11 0 12 13 12 1900 19 48

15 1918 19 16 3818 19 WANT PIUTURE EYZNJT

3800

1.8

13

14

Figure 6

18

13

18

10

19

19

```
16*[1:1]:[[[[[]]]]]
21×4444444444444444×*5555555555555555***6666666666666666*...
?**44444444444444444444.
```

1

Figure 7

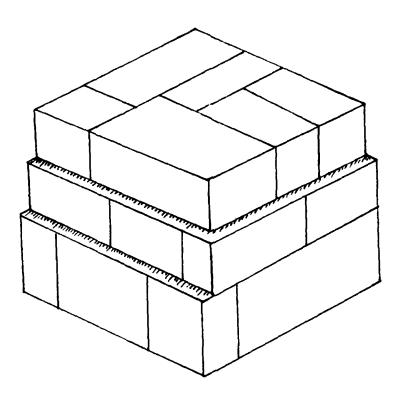


Figure 8a

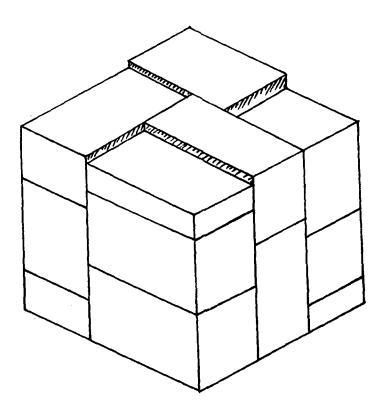


Figure 8b

C. Attempt to select—set of the lover boxer without record to common dimensions. The load the pallet. Since the top of the boxes grob ble will not be the came moreht, it will be negative to define the tops of each of the boxes (or several boxes of common beight) as "pallets" and proceed to load boxes on the tops of the "collets". This is easily done within the structure of the UYNALOAU routine since both length and width are user input and "L" shaped pallets can be specified. Combinations of this and the previous approach can be used to complete loading of the pallet.

3. Select sets of boxes using the STACK BOXES PRIOR TO PALLER LOADING option. Then use the Pallet Loading Procedure to Position the Stacks on the Pallet (figure 8b).

Pallet Loading Algorithm

The pallet loading procedure used in IFLS can be described as a combination of the principles of Dynamic Programming (D.F.) and heuristics. A serendipitous by-product of the structure of this combination is that positioning of hazardous material and considerations of center of gravity can be handled without loss within the procedure. To understand the procedure, it is useful first to consider a "best" procedure. Assume that it is desirable (no matter what the cost) to find solutions to the two-dimensional pallet loading problem which maximize the area covered on the pallet. In order to achieve this end, Dynamic Programming will be used. The following definitions will be useful.

- P = A partition dividing the pallet into two parts (see figure 9). The left-hand sub-pallet must include the origin (0,0), and the right-hand sub-pallet must include the point (L,W).
- i = The index of boxes to be loaded, i=1,...,n.
- l(i) = The length of box i.
- w(i) = The width of box i.
- S(i) = The profile (shadow) of box i. The profile is a rectangle with length l(i) and width w(i).
- N = Set of all boxes to be considered for loading (of size n).
- I = Subset of the boxes, 1,2,...,n.
- f(F,I) = The maximum area which can be covered of the lefthand sub-pallet of F using the subset of boxes I.

The Dynamic Programming equation for the pallet loading problem can be given as follows:

(1)
$$f(P,\vec{I}) = \max_{i \in I} [I(i)*\omega(i) + f(P-S(i),I-i)].$$

It should be noted that the notation 'F-S(i)' represents a partition which, in a graphical sense, is the partition F with a profile of box i removed from the right-hand edge of the left-hand sub-pallet. Typically, there could be many different realizations of 'F-S(i)' that should be considered within a dynamic optimization. There are other obvious difficulties with

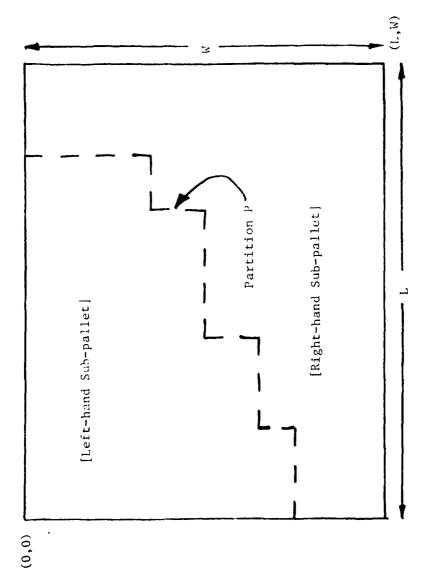


Figure 9.

Plan View of Pallet With Sample Partition P

implementing equation (1). The most obvious in that the number of mossible partition of the sallet is extremely large. The means that the state space for the D.P. will require large amounts of computer storage. It also means that the computer time required to solve the D.P. likely would be well beyond some sensible limit for real world applications.

One approach to developing a more efficient procedure in to limit in some way the form of the possible partitions of the pallet. In the present case, partitions of the pallet have been limited to rectangles (figure 10a). In order to specify the Dynamic Program resulting from the rectangular partitions, the following additional definitions are needed.

x = Two-dimensional index specifying a rectangular
partition (figure 6a).

J = Subset of the boxes, 1,2,...,n.

s(x,y,I) = The maximum area which can be covered of the lefthand sub-pallet of x,y using the subset of boxes I.

h(x,y,x',y',I) = The maximum area which can be covered of the left-hand sub-pallet of x',y' less the left-hand sub-pallet of x,y using boxes from the set I (not all elements of I necessarily are used) see figure 10b).

The dynamic programming equation for the pallet loading problem (limited to rectangular partitions) can be given as follows:

The implementation of equation (2) also has its difficulties. The function h(x,y,x',y',I) itself requires an optimization in order to pack the L-shaped area common to the partition x',y', but not common to the partition x,y' (i.e., the cross-hatched area in figure 10b). The state-space is still too large to deal with on a practical basis.

In order to limit the size of the state space, it was decided to carry only one partial solution (s(x,y,I)) for each partition x,y. The obvious choice is to carry

max [s(x,y,l)].

With this limitation on the state space, the implementation of equation (2) is relatively straightforward. It is necessary, however, to specify several important details first:

1. an optimization structure for h(x,y,x',y',I); and

2. bounding rules for the elimination of partial solutions. The optimization for h(x,y,x',y',I) is done simply by breaking the L-shaped area into two rectangular areas (figure 11) and filling each area useing a simple D.P. procedure. The following definition is useful.

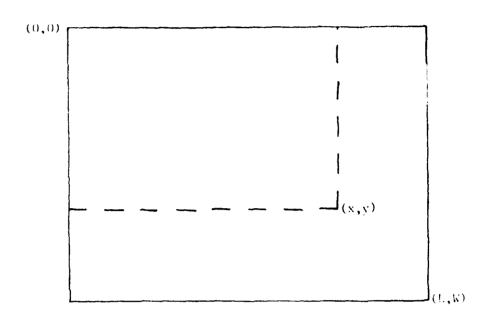


Figure 10a Pallet With Rectangular Partition (x.y)

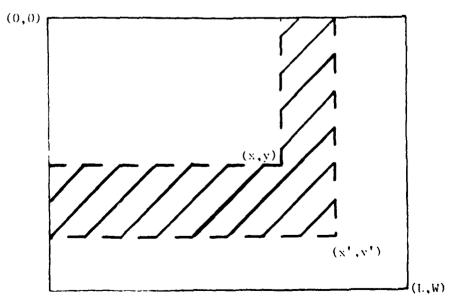
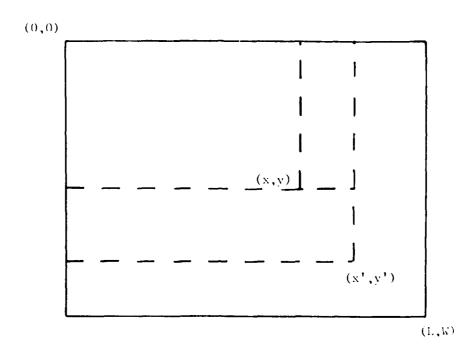


Figure 40b
Pallet With Two Rectangular Partitions
(x,y) and (x',v')



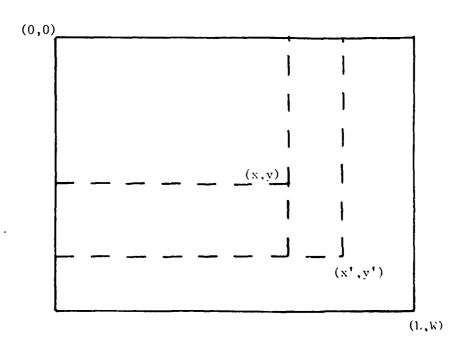


Figure 11 Two Ways To Break Up L-Shaped Area

h(1,4) The maximum rossible crace covered in a rectangular area of lemith the stacking boxes from the set 1,...i.

The Dunamic Presnamming equation for the rectangular loading problem can be diven as follows:

(β) δ(j,χ) = max [b(j 1,χ), b(j-1,χ-1(j))+1(j)*ω(j)].

The implementation of equation (3) is achieved by first turning each box in the candidate set so that its longest dimension is perpendicular to the long dimension of the rectangular area (if the longest dimension of the box is less than or equal to the short dimension of the rectangular area, that is). This insures an optimal macking of the rectangular area. The L-shaped area is broken into two rectangular areas (corridors) two different wave (figure 11) for the application of equation (3). The algorithm retains the best solution obtained.

A simple, and almost obvious, bounding rule that eliminates a great deal of the storage requirements for the state space is that a partial solution does not need to be retained for the partition x,y if there exists a partial solution for a partition x',y', $(x'\leq \pi x, y'\leq \pi y)$ such that the solution value for x',y' (max Eg(x',y',I)) is greater than or equal to the solution value for x,y (maxEg(x,y,I)).

Another effective bounding scheme is used to eliminate computation time. The following definition is useful.

The Dynamic Programming equation for the linear loading problem can be given as follows:

(4) $c(j,x) = \max \{ c(j-1,x), c(j-1,x-1(j)) + 1(j) \}$

Each box is considered by both length and width as a candidate for insertion. The function c(n,x) specifies the maximum linear coverage that is possible on the line segment [0,x] choosing from the set of boxes $1,\ldots,n$. For a pallet (rectangular sub-pallet) of size L by W, an upper bound on the maximum load (coverage) possible is c(n,L)*c(n,W). The function c(n,x) can be computed prior to the pallet loading and is easily implemented within the structure of equation 2.

Since the solution procedure results in the boxes being placed in corridors on the pallet, two serendipitous by-products occur. First, since each corridor has at least one end on the pallet perimeter, any box which contains hazardous material can be placed at the end of the corridor (as per USAF Resulations) within the structure of the solution. Second, asain since boxes can be moved within their assigned corridors, some control can be maintained over the center of aravity of the pallet within the structure of the solution.

Comments on Implementation

In order to implement the IPLS system in the USAF operating environment, there are several requirements that must be met. First, there are two elements of data that must be added to the BAMS data set; this-end-up information, and stacking capability. Observations of USAF personnel while packing pallets indicate that many items can be packed any-end-up while other items must be packed this-end-up. Adding an indicator for this to the BAMS data base is imperative to any automated packing scheme. There are many items that cannot have anything placed on top of them due to either strength or shape considerations. Further study may indicate that it is necessary for actual strength estimates to be made. However, it is our observation that the vast majority of items are either "stackable" or not. With all this in mind, an additional parmeter of the following form would be appropriate for the BAMS data:

- 0 Ans-End-Us and Stackable;
- 1 This-End-Up and Stackable;
- 2 Any-End-Up and Non-Stackable; and
- 3 This-End-UP and Non-Stackable.

Second, there is certain equipment that would be most useful. The present system was developed using a PDP-11/34 with the interface instrument being either a "dumb" CRT terminal or a hard copy terminal. This equipment reasonably duplicates the computing power (B-3500) and peripheral equipment presently available at USAF base level. However, Computer equipment has revolutionized in the last ten years. Due to its limited computing power, and limited directly addressable memory, the PDP 11/34 is marginally adequate for the Job. Test problems run on the system are solved in anywhere from a few seconds to a several minutes. Solution times are dependent on pallet size and the input data (set of boxes to be loaded). Faster CPU speeds would be desireable.

The man-machine interface for the system would be greatly facilitated with the use of a "smart" graphics CRT terminal. Three dimensional views of the pallet (with currently loaded boxes) would give the user a great deal of insight into the loading process. It is difficult to obtain that kind of insight from the present primitive two dimensional plan layout. In addition, it would be helpful to have the three-dimensional views of the pallet in hard copy output in order to facilate the actual loading personnel in implimenting the planned load.

Finally, IFLS should be considered, at this stage, as a nevelopmental system. It has not been totally "idiot proofed", nor is it completely "user friendly". Most of the shortcominas exist because hardware has not been specified for the real-life application. Once the appropriate hardware has been taigeted for the real-life application, these shortcominas can be readily corrected.

References

- Dimensional Cutting Problems, ** OPERATIONS RESEARCH, Vol. 25.
 No. 1, Jan. 1977, Pp. 30-44.
- E21 DeSha, E.L., "Area Efficient and Volume Efficient Algorithms for Loading Cargo," Masters Thesis, United States Navy Post Graduate School, Sert. 1970.
- f33 Gares, Michael, and Johnson, David, COMPUTERS AND INTRAC-TABILITY, W.H. Freeman, San Francisco, 1979.
- [43] Gilmore, P.C., and Gomory, R.E., "Multistage Cutting Stock Problems of Two and More Dimensions," OPERATIONS RESEARCH, Vol. 13, No. 1, Jan. 1965, pp.94-120.
- [5] Golden, B.L., "Approaches to the Cutting Stock Problem," ALIE TRANSACTIONS, Vol. 8, No. 2, June 1976, pp.265-272.
- [6] Hahn, s., "On the Optimal Cutting of Defective Glass Sheets." IBM N.Y. Scientific Center Report No. 320-2916, 1967.
- [7] Herz, J.C., "A Recursive Computing Procedure For Two-Dimensional Stock Cutting," IBM JOURNAL OF RESEARCH AND DEVELOPMENT, Vol. 140-1972, pp. 462-469.
- [8] Nemhauser, George L., INTRODUCTION TO DYNAMIC PROGRAMMING, John Wiley and Sons, Inc., New York, NY, 1966.
- [9] Steudel, H.J., "Generating Pallet Loading Patterns: A Special Case of the Two-Dimensional Cutting Stock Problem," MANAGEMENT SCIENCE, Vol. 25, No. 10, Oct 1979, pp. 997-1004.